

unambiguously traced to a unique origin and final recipient. Even if two participants exchange funds in an indirect way, a properly engineered path-finding method will reveal the origin and final recipient.

It is also suspected that Bitcoin does not satisfy the second property. Some researchers stated ([33, 35, 29, 31]) that a careful blockchain analysis may reveal a connection between the users of the Bitcoin network and their transactions. Although a number of methods are disputed [25], it is suspected that a lot of hidden personal information can be extracted from the public database.

Bitcoin's failure to satisfy the two properties outlined above leads us to conclude that it is not an anonymous but a pseudo-anonymous electronic cash system. Users were quick to develop solutions to circumvent this shortcoming. Two direct solutions were "laundering services" [2] and the development of distributed methods [3, 4]. Both solutions are based on the idea of mixing several public transactions and sending them through some intermediary address; which in turn suffers the drawback of requiring a trusted third party.

Recently, a more creative scheme was proposed by I. Miers et al. [28]: "Zerocoin". Zerocoin utilizes a cryptographic one-way accumulators and zero-knowledge proofs which permit users to "convert" bitcoins to zerocoins and spend them using anonymous proof of ownership instead of explicit public-key based digital signatures. However, such knowledge proofs have a constant but inconvenient size - about 30kb (based on today's Bitcoin limits), which makes the proposal impractical. Authors admit that the protocol is unlikely to ever be accepted by the majority of Bitcoin users [5].

2.2 The proof-of-work function

Bitcoin creator Satoshi Nakamoto described the majority decision making algorithm as "one-CPU-one-vote" and used a CPU-bound pricing function (double SHA-256) for his proof-of-work scheme. Since users vote for the single history of transactions order [1], the reasonableness and consistency of this process are critical conditions for the whole system.

The security of this model suffers from two drawbacks. First, it requires 51% of the network's mining power to be under the control of honest users. Secondly, the system's progress (bug fixes, security fixes, etc...) require the overwhelming majority of users to support and agree to the changes (this occurs when the users update their wallet software) [6]. Finally this same voting mechanism is also used for collective polls about implementation of some features [7].

This permits us to conjecture the properties that must be satisfied by the proof-of-work pricing function. Such function must not enable a network participant to have a *significant* advantage over another participant; it requires a parity between common hardware and high cost of custom devices. From recent examples [8], we can see that the SHA-256 function used in the Bitcoin architecture does not possess this property as mining becomes more efficient on GPUs and ASIC devices when compared to high-end CPUs.

Therefore, Bitcoin creates favourable conditions for a large gap between the voting power of participants as it violates the "one-CPU-one-vote" principle since GPU and ASIC owners possess a much larger voting power when compared with CPU owners. It is a classical example of the Pareto principle where 20% of a system's participants control more than 80% of the votes.

One could argue that such inequality is not relevant to the network's security since it is not the small number of participants controlling the majority of the votes but the honesty of these participants that matters. However, such argument is somewhat flawed since it is rather the possibility of cheap specialized hardware appearing rather than the participants' honesty which poses a threat. To demonstrate this, let us take the following example. Suppose a malevolent individual gains significant mining power by creating his own mining farm through the cheap